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6.854J / 18.415J Advanced Algorithms
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18.415/6.854 Advanced Algorithms

Problem Set 6

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1. The *betweenness problem* is defined as follows: We are given n and a set T of m triples of the elements of $\{1, \dots, n\}$. We say that an ordering π of $\{1, \dots, n\}$ satisfies a triple (i, j, k) , if j is between i and k in π . (For example, the ordering $(5, 3, 1, 2, 4)$ satisfies the triples $(5, 1, 2)$ and $(1, 3, 5)$, but not $(3, 2, 1)$). The question is to find an ordering of $\{1, \dots, n\}$ that satisfies the maximum number of triples in T .

This problem is known to be NP-hard, even if we restrict to instances for which an ordering that satisfies all the triples exist.

- (a) Use randomization to find a simple $\frac{1}{3}$ -approximation algorithm for this problem. Prove the correctness of your algorithm.
- (b) Use the method of conditional expectations to derandomize your algorithm.
- (c) Assume there is an ordering that satisfies all the triples in T . Prove that there are vectors $v_1, \dots, v_n \in \mathbb{R}^n$ such that

$$\begin{aligned} \|v_i - v_j\| &\geq 1 && \text{for all } i, j, \\ (v_i - v_j)(v_k - v_j) &\leq 0 && \text{for all } (i, j, k) \in T \end{aligned} \quad (1)$$

Show how we can find such v_1, \dots, v_n using semidefinite programming.

- (d) Give an example where the program (1) is satisfiable, but there is no ordering that satisfies all the triples in T .
 - (e) Assume that $v_1, \dots, v_n \in \mathbb{R}^n$ is a solution of the program (1). Choose r uniformly at random from $\{p \in \mathbb{R}^n : \|p\| = 1\}$, and consider the ordering obtained by sorting the elements of $\{1, \dots, n\}$ with respect to their $r^T v_i$ value. Show that in expectation this ordering satisfies at least half the constraints in T .
2. Consider the following scheduling problem. We are given n jobs that are all available at time 0 and that can be processed on any of m machines. Each job has a processing time p_j which represents the amount of time a machine (any one of them) needs to process it (without interruption). A machine can only process one job at a time. This scheduling problem is to assign each job to a machine and schedule the jobs so as to minimize $\sum_j p_j C_j$ where C_j represents the time at which the processing of job j completes. (For example, if we have 5 jobs of unit processing time and 3 machines, there are many ways of obtaining an objective function value of $1 + 1 + 1 + 2 + 2 = 7$.)
 - (a) Show that the problem is equivalent to minimizing $\sum_{i=1}^n M_i^2$ where M_i is the total amount of processing time assigned to machine i .

- (b) Let $L = \frac{1}{m} \sum_j p_j$ be the average load of any machine. Show that any optimum solution for $\sum_{i=1}^n M_i^2$ will be such that each machine i either satisfy $M_i \leq 2L$ or processes a single job j with $p_j > 2L$.
- (c) Assume that $p_j \geq \alpha L$ for some constant $\alpha > 0$ for every job j , and assume that all p_j 's can only take k different values, where k is a fixed constant. Design a polynomial-time algorithm for this case.
- (d) Assume that $p_j \geq \alpha L$ for some constant $\alpha > 0$ for every job j . Design a polynomial-time approximation scheme for this case.
- (e) Not part of the problem set and not graded, just for "fun". Can you also get a polynomial-time approximation scheme without the $p_j \geq \alpha L$ assumption?